

Classification of all associative mono- n -ary algebras with 2 elements

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ABSTRACT. We consider algebras with a single n -ary operation and a certain type of associativity. We prove that (up to isomorphism) there are exactly 5 of these associative mono- n -ary algebras with 2 elements for even $n \geq 2$ and 6 for odd $n \geq 3$. These algebras are described explicitly. It is shown that a similar result is impossible for algebras with at least 4 elements. An application concerning the assignment of a control bit to a string is given.

1. Introduction

One of the most demanding tasks of combinatorics consists in counting finite algebraic structures with certain properties or even in classifying them up to isomorphism. Various types of structures can be enumerated by the number of their elements using the ingenious methods of Pólya [6, 7]. However, there are some very difficult combinatorial problems that have not been solved until now. One of these problems is to determine the number of semigroups with k elements, where k is a positive integer. An asymptotic formula for the number of labelled semigroups with k elements was found in [3]. But a different problem consists in counting up to isomorphism. Let us reformulate this problem into the language of universal algebra. An n -ary operation μ on a set A is a function $A^n \rightarrow A$. In this article the ground set A will always be finite. (A, μ) is called *mono- n -ary algebra (with $\#A$ elements)*. Two mono- n -ary algebras (A, μ) and (B, ν) are *isomorphic* if there is a bijection f from A to B , so that for all $x_1, x_2, \dots, x_n \in A$

$$f(\mu(x_1, x_2, \dots, x_n)) = \nu(f(x_1), f(x_2), \dots, f(x_n)).$$

Hence we have asked how many associative mono-2-ary algebras with k elements exist up to isomorphism. We may consider the more general problem of determining the number of associative mono- n -ary algebras with k elements, where n and k are positive integers. But what does it mean for an algebra with a single operation to be associative?

In fact, there are many different ways to generalize associativity from binary to n -ary operations. A well-known example is superassociativity, introduced by

Date: March 17, 2004.

1991 Mathematics Subject Classification: Primary: 20N15; Secondary: 20M99; 05A15.

Key words and phrases: Associativity, mono- n -ary algebra, polyadic semigroup, n -ary operation, enumeration, combinatorics.

Menger [4]. Another example is the associativity of diagonal algebras [5]. This paper will treat only a natural kind of associativity which could be called *left-right-pushing*. Let (A, μ) be a mono- n -ary algebra. It is called *associative* (in the left-right sense) if for all $1 \leq i < n$ and $a_1, a_2, \dots, a_{2n-1} \in A$

$$\begin{aligned} & \mu(a_1, \dots, a_{i-1}, \mu(a_i, a_{i+1}, \dots, a_{n+i-1}), a_{n+i}, a_{n+i+1}, \dots, a_{2n-1}) \\ = & \mu(a_1, \dots, a_{i-1}, a_i, \mu(a_{i+1}, \dots, a_{n+i-1}, a_{n+i}), a_{n+i+1}, \dots, a_{2n-1}) \end{aligned}$$

Note that a list of the type a_{m+1}, \dots, a_m denotes an empty list.

This type of associativity implies a general associativity law that will be explained in section 2.

Left-right-pushing was introduced by Dörnte [1] who generalized the notion of groups to m -groups where the binary operation is replaced by an m -ary. These m -groups (or polyadic groups) were investigated further by Post [8]. Later the concept was widened to *polyadic semigroups* by Zupnik [9] or *associatives* as Gluskin and Shvarts [2] called them. These are other names for the associative mono- n -ary algebras we study here. Whilst Zupnik, Gluskin and Shvarts investigated the representation of certain polyadic semigroups by (associative) binary and unary operations, we are interested in counting polyadic semigroups with no restrictions on their structure, only depending on their order and arity.

We would like to complete table 1 which lists the numbers of associative mono- n -ary algebras with k elements (up to isomorphism) for some small values of n and k . The entries of the table were computed by the author using brute force algorithms.

$n \backslash k$	1	2	3	4	5	6	7	8	...
1	1	3	7	19	47	130	343	951	...
2	1	5	24	188
3	1	6
...

TABLE 1. Number of associative mono- n -ary algebras with k elements

Since it seems to be very difficult to find general formulas for the rows even for small values of n one might try to find formulas for the columns for small values of k . Surprisingly we will find a simple formula for the case $k = 2$. Moreover, we classify explicitly all associative mono- n -ary algebras with 2 elements using purely elementary methods. The occurring algebras will be introduced in 3.1. In 3.2 we prove that every associative mono- n -ary algebra with 2 elements is isomorphic to one of these algebras which will be the main subject of this article. Similar results for the columns with $k \geq 3$ could not be found, but in section 4 we will see that the entries of the columns with $k \geq 4$ are at least exponentially increasing.

These results might be interesting for some aspects of nonlinear coding theory. An example of an application is given in section 5.

2. The general associativity law

When denoting a (multi-)product of several elements of a semigroup it is not necessary to write brackets in order to indicate the order of calculation since the value of the product only depends on the elements and their total order. This is called “general associativity law”, and it may be generalized for multiproducts in associative mono- n -ary algebras.

Let us fix some terms. If (A, μ) is a mono- n -ary algebra then a *product* is a formal expression $\mu(x_1, \dots, x_n)$ with some *entries* x_1, \dots, x_n . A *bracket-expression* is defined recursively by

- (1) every element of the algebra is a bracket-expression, and
- (2) every product with bracket-expressions as entries is a bracket-expression.

A *bracket* is the beginning of a product together with the end of this product. Furthermore, a *right-tower-bracketing* is a bracket-expression where for each occurring product all entries are single elements of the algebra except for the last entry.

We can formulate now a basic result that was remarked by Dörnte [1]. Because of its importance for this article the idea of a formal proof is added.

Theorem 1. *In an associative mono- n -ary algebra the value of any bracket-expression does not depend on the structure of the brackets, only on the total order of the entries.*

Proof. Prove by induction on the number of brackets that every bracket-expression can be transformed into a right-tower-bracketing. This is done in principle by using the induction hypothesis, shifting the left-most inner bracket to the right-most position (including all complicated bracket-expressions) and using the induction hypothesis again. \square

3. The 2-elements-column

Here we consider associative mono- n -ary algebras with 2 elements. This main section is organized as follows: In 3.1 the main result is presented whereas 3.2 is devoted to its proof.

3.1. The occurring algebras. We define the following types of algebras on the set $\{0, 1\}$ with single n -ary operation μ :

type **0**:

$$\mu(x_1, x_2, \dots, x_n) := 0$$

type **A**:

$$\begin{aligned} \mu(x_1, x_2, \dots, x_n) &:= 0 && \text{if } \exists i : x_i = 0 \\ \mu(1, 1, \dots, 1) &:= 1 \end{aligned}$$

type **L**:

$$\mu(x_1, x_2, \dots, x_n) := x_1$$

type **R**:

$$\mu(x_1, x_2, \dots, x_n) := x_n$$

type **G0**:

$$\begin{aligned} \mu(x_1, x_2, \dots, x_n) &:= 0 & \text{if } \#\{i \mid x_i = 0\} \equiv 1 \pmod{2} \\ \mu(x_1, x_2, \dots, x_n) &:= 1 & \text{if } \#\{i \mid x_i = 0\} \equiv 0 \pmod{2} \end{aligned}$$

type **G1**:

$$\begin{aligned} \mu(x_1, x_2, \dots, x_n) &:= 0 & \text{if } \#\{i \mid x_i = 0\} \equiv 0 \pmod{2} \\ \mu(x_1, x_2, \dots, x_n) &:= 1 & \text{if } \#\{i \mid x_i = 0\} \equiv 1 \pmod{2} \end{aligned}$$

It is easy to see that these algebras are associative. (For type **G0** and type **G1** distinguish cases according to the parity of the number of zeros in the inner and outer brackets when verifying the associativity laws.)

In case of odd $n \geq 3$ these types are pairwise nonisomorphic. The same holds in case of even $n \geq 2$ except for type **G0** and type **G1** which are then isomorphic. In the latter case we also call type **G0** and type **G1** simply type **G**. In the trivial case $n = 1$ type **L**, type **R** and type **G0** are isomorphic to type **A**.

The types **G0** and **G1** are the only polyadic groups with 2 elements which can be seen easily [1]. However, the analogue classification of polyadic semigroups is more complicated because of the lack of invertibility.

The main result of this article will be:

Theorem 2. *Let (A, μ) be an associative mono- n -ary algebra with $\#A = 2$ and $n \geq 2$. Then (A, μ) is isomorphic to one of the types defined above.*

3.2. Proof of theorem 2. Let $\mathcal{A} = (\{0, 1\}, \mu)$ be an associative mono- n -ary algebra, $n \geq 2$. In order to simplify the notation we denote the products only with brackets, i.e.

$$(a_1 a_2 \dots a_n) := \mu(a_1, a_2, \dots, a_n).$$

The general idea of the proof is to distinguish cases according to the values of certain products. In each case \mathcal{A} will be determined only by these values.

There are four possibilities for the values of the products $(0 \dots 0)$ and $(1 \dots 1)$. The case 1,0 is treated in lemma 1, the case 0,0 in lemma 2 and lemma 3. Obviously, by exchanging 0 and 1, the case 1,1 leads to isomorphic algebras as in the case 0,0.

The case $(0 \dots 0) = 0$ and $(1 \dots 1) = 1$ is more complicated. Here we need to consider the products $(0 \dots 01)$, $(10 \dots 0)$, $(01 \dots 1)$ and $(1 \dots 10)$, too. The subcases 1,0 and 0,1 resp. 0,0 for $(0 \dots 01)$, $(10 \dots 0)$ and the subcases 0,1 and 1,0 resp. 1,1 for $(01 \dots 1)$, $(1 \dots 10)$ are examined by lemma 4 resp. lemma 5. Note that these cases are not excluding each other and may be contradictory which has no effect on our proof. The remaining case is lemma 6.

Remark. If $(1 \dots 1) = 0$ in \mathcal{A} then the value of a product depends only on the number of zeros in the product and not on the order of the elements. Indeed, whenever two products have the same number of zero entries each zero can be replaced by the product $(1 \dots 1)$ and the new expressions contain both the same number of ones and thus have the same value because of the general associativity law.

Lemma 1. *If $(0 \dots 0) = 1$ and $(1 \dots 1) = 0$ then n is odd and \mathcal{A} is isomorphic to type **G1**.*

Proof. Using the assumptions and the associativity of \mathcal{A} we obtain

$$1 = (0 \dots 0) = (\underbrace{0 \dots 0}_{n-1} (1 \dots 1)) = ((\underbrace{0 \dots 0}_{n-1} 1) \underbrace{1 \dots 1}_{n-1}).$$

Thus (in order to avoid the contradiction $1 = (1 \dots 1)$)

$$(\underbrace{0 \dots 0}_{n-1} 1) = 0. \quad (1)$$

Assume that n is even. Then, using (1) we have the contradiction

$$\begin{aligned} 0 &= (1 \dots 1) = (\underbrace{(0 \dots 0) 1 \dots (0 \dots 0) 1}_{\frac{n}{2}}) \\ &= (0 (\underbrace{0 \dots 0}_{n-1} 1) \dots 0 (\underbrace{0 \dots 0}_{n-1} 1)) = (0 \dots 0) = 1. \end{aligned}$$

So n is odd.

Claim 1.1. $(\underbrace{0 \dots 0}_{2k} \underbrace{1 \dots 1}_{n-2k}) = 0$ and $(\underbrace{0 \dots 0}_{2k+1} \underbrace{1 \dots 1}_{n-2k+1}) = 1$ for $k = 0, \dots, \frac{n-1}{2}$

We prove the claim by induction. The first assertion is clear for $k = 0$. Assume that the first assertion has been proved for all $k \leq K$ and that the second assertion has been proved for all $k < K$ and that $0 \leq 2K \leq n - 1$. Then by assumption

$$0 = (\underbrace{0 \dots 0}_{2K} \underbrace{1 \dots 1}_{n-2K}) = (\underbrace{0 \dots 0}_{2K} (0 \dots 0) \underbrace{1 \dots 1}_{n-2K-1}) = (\underbrace{0 \dots 0}_{n-1} \underbrace{0 \dots 0}_{2K+1} \underbrace{1 \dots 1}_{n-2K-1}).$$

Thus (in order to avoid the contradiction $(0 \dots 0) = 0$) we have

$$(\underbrace{0 \dots 0}_{2K+1} \underbrace{1 \dots 1}_{n-2K-1}) = 1. \quad (2)$$

Now assume that the first and the second assertion have been proved for all $k < K$ and that $0 < 2K \leq n - 1$. Then by (2)

$$1 \stackrel{(2)}{=} (\underbrace{0 \dots 0}_{2K-1} \underbrace{1 \dots 1}_{n-2K+1}) = (\underbrace{0 \dots 0}_{2K-1} (0 \dots 0) \underbrace{1 \dots 1}_{n-2K}) = (\underbrace{0 \dots 0}_{n-1} \underbrace{0 \dots 0}_{2K} \underbrace{1 \dots 1}_{n-2K}).$$

Thus (in order not to have $(0 \dots 01) = 1$ which would contradict (1))

$$(\underbrace{0 \dots 0}_{2K} \underbrace{1 \dots 1}_{n-2K}) = 1.$$

By our remark it follows immediately from claim 1.1 that a product is 0 if and only if the number of its zero entries is even. Thus \mathcal{A} is of type **G1**. \square

Lemma 2. *If $(0 \dots 0) = 0$ and $(1 \dots 1) = 0$ and*

$$\underbrace{(0 \dots 0 1)}_{n-1} = 0$$

*then \mathcal{A} is isomorphic to type **0**.*

Proof. Prove by induction that

$$\underbrace{(0 \dots 0 1 \dots 1)}_{n-k} = 0 \text{ for } k = 1, \dots, n-1.$$

The induction step ($1 \leq k \leq n-2$) is

$$\begin{aligned} 0 &= \underbrace{(0 \dots 0 1 \dots 1)}_{n-k} = \underbrace{(0 \dots 0 0 \dots 0 1)}_{n-k-1} \underbrace{1 \dots 1}_k = ((0 \dots 0) \underbrace{0 \dots 0 1 \dots 1}_{n-k-2} \underbrace{1 \dots 1}_{k+1}) \\ &= \underbrace{(0 \dots 0 1 \dots 1)}_{n-k-1} \underbrace{1 \dots 1}_{k+1}. \end{aligned}$$

Use the remark again and conclude that every product is 0. \square

Lemma 3. *If $(0 \dots 0) = 0$ and $(1 \dots 1) = 0$ and*

$$\underbrace{(0 \dots 0 1)}_{n-1} = 1$$

*then n is even and \mathcal{A} is isomorphic to type **G**.*

Proof.

$$1 = \underbrace{(0 \dots 0 1)}_{n-1} = \underbrace{(0 \dots 0 (1 \dots 1) 1)}_{n-2} = ((0 \dots 0 11) \underbrace{1 \dots 1}_{n-1})$$

In order to avoid the contradiction $(1 \dots 1) = 1$ we have

$$\underbrace{(0 \dots 0 11)}_{n-2} = 0, \tag{3}$$

$$\underbrace{(0 1 \dots 1)}_{n-1} = 1. \tag{4}$$

Claim 3.1. $\underbrace{(0 \dots 0 1 \dots 1)}_{n-2k} = 0$ for $k = 1, \dots, \lfloor \frac{n-1}{2} \rfloor$

For $k = 1$ this is (3). Assume the claim has been proved for $1 \leq k \leq \lfloor \frac{n-3}{2} \rfloor$. Then

$$\begin{aligned} 0 &= \underbrace{(0 \dots 0 1 \dots 1)}_{n-2k} \stackrel{(3)}{=} \underbrace{(0 \dots 0 (0 \dots 0 11) 1 \dots 1)}_{n-2k-1} = ((0 \dots 0) \underbrace{0 \dots 0 1 \dots 1}_{n-2k-3} \underbrace{1 \dots 1}_{2k+2}) \\ &= \underbrace{(0 \dots 0 1 \dots 1)}_{n-2k-2} \underbrace{1 \dots 1}_{2k+2}, \end{aligned}$$

i.e. the claim is true for $k+1$.

Claim 3.2. $\underbrace{(0 \dots 0 1 \dots 1)}_{2k-1} = 1$ for $k = 1, \dots, \lfloor \frac{n-1}{2} \rfloor$

For $k = 1$ this is (4). Assume the claim has been proved for $1 \leq k \leq \lfloor \frac{n-3}{2} \rfloor$. Then

$$\begin{aligned} 1 &= (\underbrace{0 \dots 0}_{2k-1} \underbrace{1 \dots 1}_{n-2k+1}) \stackrel{(4)}{=} (\underbrace{0 \dots 0}_{2k-1} \underbrace{0 \underbrace{1 \dots 1}_{n-1}}_{n-2k} \underbrace{1 \dots 1}_{n-2k}) = (\underbrace{0 \dots 0}_{2k} \underbrace{1 \dots 1}_{n-2k-1}) \\ &= (\underbrace{0 \dots 0}_{2k+1} \underbrace{1 \dots 1}_{n-2k-1}), \end{aligned}$$

i.e. the claim is true for $k + 1$.

If n is odd then the two claims contradict each other. So n is even and according to the remark at the beginning of this subsection, the premises and the two claims a product is 0 if and only if it contains an even number of zeros, i.e. \mathcal{A} is of type **G**. \square

Lemma 4. If $(0 \dots 0) = 0$ and $(1 \dots 1) = 1$ and

$$(\underbrace{0 \dots 0 1}_{n-1}) = 1 \text{ and } (1 \underbrace{0 \dots 0}_{n-1}) = 0$$

then \mathcal{A} is isomorphic to type **R**.

Proof. For all $x_1, x_2, \dots, x_{n-1} \in \{0, 1\}$ we have

$$(x_1 \dots x_{n-1} 0) = (x_1 \dots x_{n-1} (\underbrace{1 \underbrace{0 \dots 0}_{n-1}}_{n-1})) = ((x_1 \dots x_{n-1} 1) \underbrace{0 \dots 0}_{n-1}) = 0 \quad (5)$$

and

$$\begin{aligned} (x_1 \dots x_{n-1} 1) &= (x_1 \dots x_{n-1} (\underbrace{0 \dots 0 1}_{n-1})) = ((x_1 \dots x_{n-1} 0) \underbrace{0 \dots 0 1}_{n-2}) \\ &\stackrel{(5)}{=} (\underbrace{0 \dots 0 1}_{n-1}) = 1. \end{aligned}$$

Thus \mathcal{A} is isomorphic to type **R**. \square

If we had instead of $(0 \dots 0 1) = 1$ and $(1 0 \dots 0) = 0$ the alternative conditions $(1 \dots 1 0) = 0$ and $(0 1 \dots 1) = 1$ in lemma 4 we also would obtain type **R** since exchanging 0 and 1 does not affect this type. On the other hand, if we exchanged the order of the entries of all products, i.e. if we had $(0 \dots 0 1) = 0$ and $(1 0 \dots 0) = 1$ (or $(1 \dots 1 0) = 1$ and $(0 1 \dots 1) = 0$) \mathcal{A} would be isomorphic to type **L**.

By a similar argument one could change the conditions $(0 \dots 0 1) = 0$ and $(1 0 \dots 0) = 0$ in the next lemma into $(1 \dots 1 0) = 1$ and $(0 1 \dots 1) = 1$.

Lemma 5. If $(0 \dots 0) = 0$ and $(1 \dots 1) = 1$ and

$$(\underbrace{0 \dots 0 1}_{n-1}) = 0 \text{ and } (1 \underbrace{0 \dots 0}_{n-1}) = 0$$

then \mathcal{A} is isomorphic to type **A**.

Proof. For all $x_1, x_2, \dots, x_n \in \{0, 1\}$ we have

$$(0 x_2 \dots x_n) = ((\underbrace{0 \dots 0 1}_{n-1}) x_2 \dots x_n) = (\underbrace{0 \dots 0}_{n-1} (1 x_2 \dots x_n)) = 0, \quad (6)$$

$$(x_1 \dots x_{n-1} 0) = (x_1 \dots x_{n-1} (\underbrace{1 \underbrace{0 \dots 0}_{n-1}}_{n-1})) = ((x_1 \dots x_{n-1} 1) \underbrace{0 \dots 0}_{n-1}) = 0, \quad (7)$$

and for arbitrary $1 \leq k \leq n-2$

$$\begin{aligned} (\underbrace{1 \dots 1}_k 0 x_{k+2} \dots x_{n-1} 1) &= (\underbrace{1 \dots 1}_k (\underbrace{1 0 \dots 0}_{n-1}) x_{k+2} \dots x_{n-1} 1) \\ &= (\underbrace{1 \dots 1}_{k+1} \underbrace{0 \dots 0}_{n-k-2} (\underbrace{0 \dots 0}_{k+1} x_{k+2} \dots x_{n-1} 1)) \stackrel{(6)}{=} (\underbrace{1 \dots 1}_{k+1} \underbrace{0 \dots 0}_{n-k-1}) \stackrel{(7)}{=} 0. \end{aligned}$$

Thus \mathcal{A} is isomorphic to type **A**. □

Lemma 6. *If $(0 \dots 0) = 0$ and $(1 \dots 1) = 1$ and*

$$(\underbrace{0 \dots 0}_{n-1} 1) = 1 \text{ and } (1 \underbrace{0 \dots 0}_{n-1}) = 1 \text{ and } (\underbrace{1 \dots 1}_{n-1} 0) = 0 \text{ and } (0 \underbrace{1 \dots 1}_{n-1}) = 0$$

*then n is odd and \mathcal{A} is isomorphic to type **G0**.*

Proof. We need some preparations before realizing that for all types of entries of a product the values are fixed by the premises. Note that we may not use the remark, since the algebra restricted on the main diagonal is the identical combination.

For $n = 2$ the premises are contradictory. So assume that $n \geq 3$.

Claim 6.1. $(\underbrace{0 \dots 0}_{n-2} 11) = 0$

Assume that the assertion is not true, i.e. that

$$(\underbrace{0 \dots 0}_{n-2} 11) = 1.$$

Then we obtain the contradiction

$$\begin{aligned} 0 &= (\underbrace{0 1 \dots 1}_{n-1}) = (0 (\underbrace{0 \dots 0}_{n-2} 11) \underbrace{1 \dots 1}_{n-2}) = ((\underbrace{0 \dots 0}_{n-1} 1) \underbrace{1 \dots 1}_{n-1}) \\ &= (1 \dots 1) = 1. \end{aligned}$$

Claim 6.2. $(\underbrace{0 \dots 0}_{n-2k+1} \underbrace{1 \dots 1}_{2k-1}) = 1$ for $k = 1, \dots, \lfloor \frac{n}{2} \rfloor$

The proof is by induction. The case $k = 1$ is given by a premise. Now assume that $1 \leq k \leq \lfloor \frac{n-2}{2} \rfloor$ and that the assertion is true for k . Then by assumption and claim 6.1

$$\begin{aligned} 1 &= (\underbrace{0 \dots 0}_{n-2k+1} \underbrace{1 \dots 1}_{2k-1}) = (\underbrace{0 \dots 0}_{n-2k} (\underbrace{0 \dots 0}_{n-2} 11) \underbrace{1 \dots 1}_{2k-1}) = ((\underbrace{0 \dots 0}_{n-2k-2} \underbrace{0 \dots 0}_{2k-1}) \underbrace{1 \dots 1}_{2k+1}) \\ &= (\underbrace{0 \dots 0}_{n-2k-1} \underbrace{1 \dots 1}_{2k+1}), \end{aligned}$$

thus the assertion is true for $k+1$.

Claim 6.3. n is odd.

Indeed, if n was even we would have

$$(0 \underbrace{1 \dots 1}_{n-1}) = 0$$

by prerequisite and

$$(0 \underbrace{1 \dots 1}_{n-1}) = 1$$

by claim 6.2 which is a contradiction.

We are now ready to prepare the main argument in the proof of lemma 6. If we write $(\dots x_1 x_2 \dots x_k \dots)$ the dots at the beginning and at the end denote an arbitrary number of arbitrary entries, so that the total number of entries is n . Let $2 \leq b \leq n-1$ be even. By claim 6.2 we obtain

$$\begin{aligned} (\dots \underbrace{0 \dots 0}_b \dots) &= (\dots \underbrace{0 \dots 0}_{b-1} \underbrace{0 1 \dots 1}_{n-1} \dots) = (\dots (\underbrace{0 \dots 0}_b \underbrace{1 \dots 1}_{n-b}) \underbrace{1 \dots 1}_{b-1} \dots) \\ &= (\dots \underbrace{1 \dots 1}_b \dots). \end{aligned} \quad (8)$$

According to (8) we may shift even numbers of zeros (resp. ones) as ones (resp. zeros) to the left without changing the value of a product, so that every product can be written in an equivalent *normal form*:

$$(\underbrace{0 \dots 0}_{n-2k} \underbrace{10 \dots 10}_k), \quad k = 0, 1, 2, \dots, \frac{n-1}{2} \quad (9)$$

$$(\underbrace{1 \dots 1}_{n-2k} \underbrace{01 \dots 01}_k), \quad k = 0, 1, 2, \dots, \frac{n-1}{2} \quad (10)$$

The next claims determine the values of these normal forms.

Claim 6.4. $(\underbrace{1 \dots 1}_{n-2a} \underbrace{01 \dots 01}_a) = 0$ for $a \in \{1, \dots, \frac{n-1}{2}\}$ odd and
 $(\underbrace{1 \dots 1}_{n-2b} \underbrace{01 \dots 01}_b) = 1$ for $b \in \{1, \dots, \frac{n-1}{2}\}$ even

Claim 6.5. $(\underbrace{0 \dots 0}_{n-2a} \underbrace{10 \dots 10}_a) = 1$ for $a \in \{1, \dots, \frac{n-1}{2}\}$ odd and
 $(\underbrace{0 \dots 0}_{n-2b} \underbrace{10 \dots 10}_b) = 0$ for $b \in \{1, \dots, \frac{n-1}{2}\}$ even

By symmetry arguments it is sufficient to prove claim 6.4. We use the same kind of double-step-induction as in the proof of claim 1.1.

Suppose

$$(\underbrace{1 \dots 1}_{n-2} 01) = 1.$$

Then we obtain the contradiction

$$\begin{aligned} 0 &= (\underbrace{1 \dots 1}_{n-1} 0) = (\underbrace{1 \dots 1}_{n-2} \underbrace{1 \dots 1}_{n-2} 01) (\underbrace{0 1 \dots 1}_{n-1}) \\ &= (\underbrace{1 \dots 1}_{n-3} \underbrace{1 \dots 1}_{n-1} 0) (\underbrace{10 \dots 1}_{n-2} \underbrace{1 \dots 1}_{n-2} 1) \stackrel{(8)}{=} (\underbrace{1 \dots 1}_{n-3} \underbrace{1 \dots 1}_{n-1} 0) (\underbrace{1 \dots 1}_{n-2} 01) 1 \\ &= (\underbrace{1 \dots 1}_{n-3} \underbrace{1 \dots 1}_{n-1} 0) 11 = (\underbrace{1 \dots 1}_{n-2} \underbrace{1 \dots 1}_{n-2} 01) 1 = (1 \dots 1) = 1. \end{aligned}$$

Thus

$$\underbrace{(1 \dots 1 0 1)}_{n-2} = 0.$$

Now let $B \geq 2$ be even and the assertion be true for all $a, b < B$. Then

$$\begin{aligned} \underbrace{(1 \dots 1 0 1 \dots 0 1)}_{n-2B} &= \underbrace{(1 \dots 1 0 1 \dots 0 1)}_{n-2B} \underbrace{(0 1 \dots 1)}_{n-1} 1 \\ &= \underbrace{(1 \dots 1 0)}_{n-2B} \underbrace{(1 0 \dots 1 0)}_{B-1} \underbrace{(1 \dots 1)}_{n-2B+2} \underbrace{(1 \dots 1)}_{2B-2} \stackrel{(8)}{=} \underbrace{(1 \dots 1 0)}_{n-2B} \underbrace{(1 \dots 1)}_{n-2B+2} \underbrace{(0 1 \dots 0 1)}_{B-1} \underbrace{(1 \dots 1)}_{2B-2} \\ &= \underbrace{(1 \dots 1 0 0)}_{n-2B} \underbrace{(1 \dots 1)}_{2B-2} \stackrel{(8)}{=} (1 \dots 1) = 1. \end{aligned}$$

Finally let $A \geq 3$ be odd and the assertion be true for all $a, b < A$. Then

$$\begin{aligned} \underbrace{(1 \dots 1 0 1 \dots 0 1)}_{n-2A} &= \underbrace{(1 \dots 1 0 1 \dots 0 1)}_{n-2A} \underbrace{(0 1 \dots 1)}_{n-1} 1 \\ &= \underbrace{(1 \dots 1 0)}_{n-2A} \underbrace{(1 0 \dots 1 0)}_{A-1} \underbrace{(1 \dots 1)}_{n-2A+2} \underbrace{(1 \dots 1)}_{2A-2} \stackrel{(8)}{=} \underbrace{(1 \dots 1 0)}_{n-2A} \underbrace{(1 \dots 1)}_{n-2A+2} \underbrace{(0 1 \dots 0 1)}_{A-1} \underbrace{(1 \dots 1)}_{2A-2} \\ &= \underbrace{(1 \dots 1 0 1 \dots 1)}_{n-2A} \stackrel{(8)}{=} \underbrace{(1 \dots 1 0 1)}_{n-2} = 0. \end{aligned}$$

This completes the proof of claim 6.4. Now the value of every product is determined. Observe that these values are 0 if and only if the number of zero entries is odd, so \mathcal{A} is isomorphic to type **G0**. This completes the proof of the lemma. \square

As described in the beginning of this subsection and by the notice between lemma 4 and lemma 5 the case distinctions are complete which proves theorem 2.

4. The columns with several elements

A classification of the associative mono- n -ary algebras with 3 elements has not been completed successfully until now. It would be an interesting question whether the number of these algebras is bounded like in the case of 2 elements.

If we consider the case of 4 or more elements the numbers in the column of table 1 are not bounded but at least exponentially growing. Indeed every mono- n -ary algebra with at least four elements 0, 1, 2, 3 of the following type is associative:

- If there is a 0 or an 1 entry then the product is 0.
- Otherwise the product is 0 or 1.

There are at least 2^n such algebras, and at least 2^{n-1} pairwise nonisomorphic among.

The question arises whether or not a classification of the associative mono- n -ary algebras with more than 3 elements in finitely many series of algebras exists.

5. An application: control bits

As an application of theorem 2 we study **recursive allocation of control bits** for a given word of length $1 + \ell(n - 1)$ over the alphabet $\{0, 1\}$ **taking into consideration equal inner structures** of length n .

To be precise we want to assign a single control bit i to a very long string $a_1 a_2 \dots a_{1+\ell(n-1)}$, ($n \geq 2, \ell \gg 1$). Regarding our available memory space we do not want to use a map $\{0, 1\}^{1+\ell(n-1)} \longrightarrow \{0, 1\}$, but only a map

$$\mu : \{0, 1\}^n \longrightarrow \{0, 1\}$$

which we apply recursively on substrings, replacing these substrings by the value of the mapping. For reasons of efficiency it might be advantageous not to choose the substrings in canonical order, rather in a randomlike order. This may be the case if it is cheaper to find and calculate two identical substrings than to calculate two different substrings. In order to have for each possible word a well defined control bit i (that is independent of the choice of the substrings) we must claim that μ is associative (in the left-right sense). So only associative mono- n -ary algebras $(\{0, 1\}, \mu)$ are a solution to this (very special) control bit allocation problem.

From theorem 2 we know that there are only 8 of these algebras. (Note that type **A** and type **0** appear twice since their automorphism group is trivial.) The control bits resulting from type **0**, type **A**, type **L** and type **R** are somewhat pathological, e.g. for type **L** it is simply the repetition of the first bit of the string. In this context the commonly used control bits are those depending on the parity of the stringsum and stringlength produced by the polyadic groups **G0** and **G1**, and by theorem 2 there is no sensible alternative.

The situation changes for an alphabet with $k \geq 3$ letters. In this case we should speak rather of control characters than of control bits. Even there the majority of associative mono- n -ary algebras which are not polyadic groups will lead to pathological control characters, see the algebras in 4. But there exist associative mono- n -ary algebras that are not polyadic groups with interesting associated control characters, e.g. for $n = 2$ and $k = 3$ consider the addition of a neutral element to the group $\mathbb{Z}/2\mathbb{Z}$:

μ	0	1	2
0	0	1	0
1	1	0	1
2	0	1	2

Although the control character produced by this semigroup contains hardly any information about the number of the letters 2 and 0 in the string it provides a lot of information about letter 1, namely the parity of its occurrence.

Acknowledgements

The author would like to thank A. Romanowska, P. Burmeister, U. Faigle, B. Fuchs and B. Wienand for some hints and comments.

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